# Improved Algorithm for Solving the Maximal Clique Problem with DNA Computing 

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#### Abstract

In this paper, we improve an algorithm proposed by Q. Ouyang to solve a maximal clique problem. While we can determine the size of the largest clique by sorting the length of the DNA strands in Q. Ouyang proposal, we cannot determine the exact vertices involved in the maximal clique. Thus in our improved algorithm, the exact vertices for each clique can be derived straightforwardly from the read-out in the gel electrophoresis process.


Keywords: DNA computing, the maximal clique problem, the parallel overlap assembly

## I. INTRODUCTION

DNA computing which uses bio-molecules as a basis for computation is thought one day to be able to outperform electronic computers in computing complex combinatorial problems such as NP hard. Since then, many proposals to solve NP hard problems with DNA computing have been published including an algorithm proposed by Q. Ouyang [1] for solving a maximal clique problem (MCP) in 1997. The MCP is encoded in DNA sequences by assigning a binary number to each possible clique in the graph as the data structure. A bit set to 1 represents a vertex in a clique and a bit set to 0 represents a vertex out of the clique. Therefore for a graph with N vertices, the complete set of possible cliques is ensemble into an N digit binary numbers. The data structure is then designed in the form of double stranded DNA (dsDNA) and restriction enzyme sites are embedded for each bit value $=1$. These restriction enzyme sites provide a "cutting" reaction to the enzymes and were not amplified during PCR. Reading the size of the largest clique is then based on the shortest length of the elongated DNA strands.

In our work, we improve the algorithm proposed by Q. Ouyang for better qualitative results of the computation. While we can determine the size of the largest clique by sorting the length of the DNA strands in Q . Ouyang proposal, we cannot determine the exact vertices involved in the maximal clique. Thus in our improved algorithm, the exact vertices for each clique can be derived straightforwardly from the read-out in the gel electrophoresis process.

## II. THE MAXIMAL CLIQUE PROBLEM

A complete graph is a simple graph in which every pair of distinct vertices is connected by an edge. The MCP is to find a maximum size complete sub-graph, a maximal clique, from a given graph. The MCP is proven
to be an NP-hard problem, and is difficult to solve efficiently.

## III. ALGORITHM

We discuss an algorithm to solve the MCP by showing the graph example in Fig.1. Consider a graph with eight vertices, sixteen edges and three maximal cliques. Fig. 3 shows an adjacency matrix of the given graph. The complementary graph (Fig.4) contains all edges missing in the original graph.


Fig.1. The graph of our example

| Showing the connection of the node $\mathrm{H}_{1}$ |
| :--- |
|  |
|  |$N_{1} N_{2} N_{3}$

Fig.2. Adjacency matrix for the graph


Fig.3. The complementary graph
The new algorithm consists of the following six steps.

## Step 1 Cording (DNA sequence Generator)

First, we construct the sequence for value $N_{i}$ and position $E_{i}$ as described in Fig.1. There are eight value sections ( $N_{1}-N_{8}$ ) and its position ( $E_{1}-E_{9}$ ). The position sequences are used for connecting each node $N_{i}(i=0, \ldots, 8)$. The eight value sections are set sandwiched between nine positions. Each combined oligonucleotide consists of two position motifs, $E_{i+1} N_{i} E_{i}$ for odd $i, \overline{E_{i+1} N_{i} E_{i}}$ for even $i$, where the bar shows the complementary sequence and the value of $N_{i}$ can be either 1 or 0 . In this experiment, we designed the DNA sequences in the form of ssDNA. We set $E_{i}=$ $10 \mathrm{bp}, N_{i}=0 \mathrm{bp}$ (if the value $=1$ ) and $N_{i}=15 \mathrm{bp}$ (if the value $=0$ ). The sequence of vertices are constructed from restriction enzyme (RE) sites for each bit value $=1$.

## Step 2 Merge (Parallel Overlap Assembly)

We produce a pool containing all possible solution with eight tubes T1-T8. Following Fig.2, we add combined ENE DNA sequences including nodes $N_{i}^{1}$ showing value 1 and all nodes $N_{i}^{0}$ showing value 0 into each tube. By performing POA, the DNA sequences for a position are annealed to the complementary sequence for the other DNA sequences and the combinations of $N_{1} N_{2} N_{3} N_{4} N_{5} N_{6}$, a pool containing possible solution, is generated.

## Step 3 Cutting (Restriction Enzyme)

Contents of T1-T8 are digested with each RE (RE1RE8) to eliminate illegally connected edges (Fig. 3) in each test tube.

## Step 4 Amplify (Polymerase Chain Reaction)

Amplify only strands with $E_{0}$ and $\overline{E_{9}}$ for each tube by supplying primers $E_{0}$ and $\overline{E_{9}}$.

## Step 5 Extract (Gel Electrophoresis)

The clique of the largest size is represented by the shortest DNA length. From gel electrophoresis process, extract only the shortest band 150bp (clique size $=4$ ) for T1-T8 and separate into tubes T1'-T8'. Add RE1RE8 into tubes T1'-T8' and observe the band lengths.

## Step 6 Identify

The experiment result is observed by performing gel electrophoresis. Each vertex involved in individual
cliques can be identified separately by the clustering of lengths for the vertices. From the result, you can find the vertices $(1,2,3,4),(2,3,4,5)$, and $(4,5,6,7)$ form the maximal cliques with size $=4$.


Fig.4. Gel Electrophoresis

## VI. CONCLUSION

The experiment was carried out using basic biomolecular tools as used such as parallel overlap assembly, polymerase chain reaction (PCR) and gel electrophoresis process. A major difference in our proposal is the inclusion of Step 6. In the original algorithm by Q. Ouyang, separate vertices involved in a clique were confirmed using molecular cloning technique which is time and resources inefficient. Using our algorithm the size of each clique can be determined by its DNA strand length and each vertex involved in individual cliques can be identified separately.

## REFERENCES

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