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Cooperating Systems of Three-Dimensional Finite Automata

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Abstract

In 1967, M.Blum and C.Hewitt first proposed two-dimensional automata as a computational model of two-dimensional pattern processing, and investigated their pattern recognition abilities. Since then, a lot of researchers in this field have been investigating many properties about automata on a two-dimensional tape. On the other hand, the question of whether processing three-dimensional digital patterns is much more difficult than two-dimensional ones is of great interest from the theoretical and practical standpoints. Thus, the study of three-dimensional automata as a computasional model of three-dimensional pattern processing has been meaningful. This paper introduces a cooperating system of three-dimensional finite automata as one model of three-dimensional automata. A cooperating system of three-dimensional finite automata consists of a finite number of three-dimensional finite automata and a three-dimensional input tape where these finite automata work independently (in parallel). Those finite automata whose input heads scan the same cell of the input tape can communicate with each other, that is, every finite automaton is allowed to know the internal states of other finite automata on the same cell it is scanning at the moment. In this paper, we mainly investigate several accepting powers of a cooperating system of five-way three-dimensional finite automata. The five-way three-dimensional finite automaton is a three-dimensional finite automaton whose input head can move east,west,south,north,or down,but not up on a three-dimensional input tape.

Key Words **:** computational complexity, cooperating system, finite automaton, multihead, three-dimension

1 Introduction

A cooperating system of three-dimensional finite automata (CS-3-FA) [2,3,4] consists of a finite number of three-dimensional finite automata and a three-dimensional input tape where these finite automata work independently (in parallel). Those finite automata whose input heads scan the same cell of the input tape can communicate with each other, that is, every finite automaton is allowed to know the internal states of other finite automata on the same cell it is scanning at the moment.

In this paper, we propose a cooperating system of five-way three-dimensional finite automata (CS-FV3-FA) which is a restricted version of CS-3-FA's, and mainly investigate its several properties as three-dimensional language acceptors. The five-way three-dimensional finite automaton [7] is a three-dimensional finite automaton [1] whose input head can move east,west,south,north,or down, but not up.

The paper has six sections in addition to this Introduction. Section 2 contains some definitions and notions. Section 3 investigates a relationship between five-way three-dimensional simple multihead finite automata (FV3-SPMHFA's) and CS-FV3-FA's. It is shown that FV3-SPMHFA's and CS-FV3-FA's are equivalent in accepting power if the input tapes are restricted to cubic ones. Section 4 investigates the difference between the accepting powers of CS-FV3-FA's and CS-3-FA's, and shows that CS-FV3-FA's are less powerful than CS-3-FA's. Section 5 investigates the difference between the accepting powers of deterministic and nondeterministic CS-FV3-FA's, and shows that deterministic CS-FV3-FA's are less powerful than nondeterministic CS-FV3-FA's. Section 6 concludes by giving some open problems. In this paper only cubic input tapes are considered.

2 Preliminaries

Definition 2.1. Let $\sum$ be a finite set of symbols. A *three-dimensional tape* over $\sum$ is a three-dimensional rectangular array of elements of $\sum$. The set of all three-dimensional tapes over $\sum$ is denoted by $\sum^{(3)}$. Given a tape $x \in \sum^{(3)}$, for each integer $j (1 \leq j \leq 3)$, we let $l_j(x)$ be the length of x along the jth axis. The set of all $x \in \sum^{(3)}$ with $l_1(x) = n_1$, $l_2(x) = n_2$, and $l_3(x) = n_3$ is denoted by $\sum^{(n_1,n_2,n_3)}$. When $1 \leq i_j \leq l_j(x)$ for each $j (1 \leq j \leq 3)$, let $x(i_1, i_2, i_3)$ denote the symbol in x with coordinates (i_1, i_2, i_3). Furthermore, we define

$$x[(i_1, i_2, i_3), (i'_1, i'_2, i'_3)],$$

when $1 \leq i_j \leq i'_j \leq l_j(x)$ for each integer $j(1 \leq j \leq 3)$, as the three-dimensional input tape y satisfying the follwing conditions:

(i) for each $j(1 \leq j \leq 3)$, $l_j(y) = i'_j - i_j + 1$;

(ii) for each $r_1, r_2, r_3(1 \leq r_1 \leq l_1(y), 1 \leq r_2 \leq l_2(y), 1 \leq r_3 \leq l_3(y)), y(r_1, r_2, r_3) = x(r_1 + i_1 - 1, r_2 + i_2 - 1, r_3 + i_3 - 1)$. (We call $x[(i_1, i_2, i_3), (i'_1, i'_2, i'_3)]$ the $[(i_1, i_2, i_3), (i'_1, i'_2, i'_3)]$-*segment* of x.)

We recall a *five-way three-dimentional simple k-head finite automaton* (FV3-SPk-HFA)[5,6]. An FV3-SPk-HFA M is a finite automaton with k read-only input heads operating on a three-dimentional input tape surrounded by boundary symbols ♯'s. The only one head(called the '*reading*' head) of M is capable of distinguishing the symbols in the input alphabet, and the other heads(called *'counting'* heads) of M can only detect whether they are on the boundary symbols or a symbol in the input alphabet. When an input tape x is a presented to M, M determines the next state of the finite control, the next move direction (east,west,south,north,down,or no move) of each input head, depending on the present state of the finite contorol, the symbol read by the reading head, and on whether or not the symbol read by each counting head is boundary symbol. We say that M *accepts* x if M, when started in its initial state with all its input heads on $x(1, 1, 1)$, eventually halts in an accepting state with all its heads on the bottom boundary symbols of x. As usual, we define nondeterministic and dterministic FV3-SPk-HFA's.

A *five-way three-dimentional sensing simple k-head finite automaton*(FV3-SNSPk-HFA) is the same device as a FV3-SPk-HFA except that the former can detect coincidence of the input heads.

We denote a deterministic(nondeterministic) FV3-SPk-HFA by FV3-SPk-HDFA(FV3-SPk-HNFA), and denote a deterministic (nondeterministic)FV3-SNSPk-HFA by FV3-SNSPk-HDFA(FV3-SNSPk-HNFA).

We now give formal definition of a *cooperating system of k three-dimensional deterministic finite automata* (CS-3-DFA(k)) as an acceptor.

Definition 2.2. A CS-3-DFA(k) is a k-tuple $M = (\mathrm{FA}_1, \mathrm{FA}_2, \ldots, \mathrm{FA}_k)$, $k \geq 1$, such that for each $1 \leq i \leq k$,

$$\mathrm{FA}_i = (\textstyle\sum, Q_i, X_i, \delta_i, q_0 i, F_i, \phi, \sharp),$$

where

1. $\sum$ is a finite set of *input symbols*.
2. Q_i is a finite set of *states*.
3. $X_i = (Q_1 \cup \{\phi\}) \times \cdots \times (Q_{i-1} \cup \{\phi\}) \times (Q_{i+1} \cup \{\phi\}) \times \cdots \times (Q_k \cup \{\phi\})$, where '$\phi$' is a special state not in $(Q_1 \cup Q_2 \cup \cdots \cup Q_k)$.
4. $\delta_i : (\sum \cup \{\sharp\}) \times X_i \times Q_i \rightarrow Q_i \times \{$east$(= (0, +1, 0))$, west$(= (0, -1, 0))$,south$(= (+1, 0, 0))$, *north*(=(-1,0,0)),up$(= (0, 0, -1))$, *down*(=(0,0,+1))}, no move$(= (0, 0, 0))\}$ is the *next move function*, where '♯' is the *boundary symbol* not in $\sum$.
5. $q_{0i} \in Q_i$ is the *initial state* of FA_i.
6. $F_i \subseteq Q_i$ is the set of *accepting states* of FA_i.

Every automaton of M independently (in parallel) works step by step on the same three-dimentional tape x over $\sum$ surrounded by boundary symbols ♯'s. Each step is assumed to require exactly one time for its completion. For each $i(1 \leq i \leq k)$, let q_i be the state of FA_i at time 't'. Then each FA_i, enters the next state 'p_i' at time '$t + 1$' according to the function

$$\delta_i(x(\alpha, \beta, \gamma), (q'_1, \ldots, q'_{i-1}, q'_{i+1}, \ldots, q'_k), q_i) = (p_i, (d_1, d_2, d_3)),$$

where $x(\alpha, \beta, \gamma)$ is the symbol read by the input head of FA_i at time 't' and for each $j \in \{1, \ldots, i-1, i+1, \ldots, k\}$,

$$q'_j = \begin{cases} \mathrm{q}_j \in Q_j & \text{if the input heads of } \mathrm{FA}_i \text{ and } \mathrm{FA}_j\text{, are on the same input position at the moment '}t\text{';} \\ \phi & \text{otherwise,} \end{cases}$$

and moves 1st input head to $x(\alpha + d_1, \beta + d_2, \gamma + d_3)$ at time '$t + 1$'. We assume that the input head of FA_i never falls off the tape beyond boundary symbols.

When an input tape $x \in \sum^{(3)}$ is presented to M, we say that M *accepts* the tape x if each automaton of M, when started in its initial state with its input head on $x(1, 1, 1)$, eventually enters an accepting state with its input head on one of the bottom boundary symbols.

We next introduce a *cooperating system of k five-way three-dimensional deteministic finite automata* (CS-FV3-DFA(k)), with which we are mainly concerned in this paper.

Definition 2.3. A CS-FV3-DFA(k) is a CS-3-DFA(k) $M = (\mathrm{FA}_1, \mathrm{FA}_2, \ldots, \mathrm{FA}_k)$ such that the input head of each FA_i can only move east, west, south, north, or down, but not up.

To give the formal definition of a *cooperating system of k three-dimensional nondeterministic finite automata* (CS-3-NFA(k)) and a *cooperating system of k five-way three-dimensional nondeterministic finite* *automata* (CS-FV3-NFA(k)) is left to the reader.

For each $X \in$ {FV3-SPk-HDFA,FV3-SPk-HNFA,FV3-SNSPk-HDFA,FV3-SNSPk-HNFA,CS-3-DFA(k),CS-3-NFA(k),CS-FV3-DFA(k),CS-FV3-NFA(k)}, by X^c we denote an X whose input tapes are restricted to cubic ones; by $\mathcal{L}[X](\mathcal{L}([X^c])$ we denote the class of sets of input tapes accepted by X's(X^c's). We will focuse our attention on the acceptors whose input tapes are restricted to cubic ones.

3 Relationship between FV3-SPMHFA's and CS-FV3-FA's

In this section, we establish a relation between the accepting powers of five-way three-dimensional simple multihead finite automata and cooperating systems of five-way three-dimensional finite sutomata over cubic input tapes. This result will be used in the latter sections.

Lemma 3.1. *For any* $k \geq 1$ *and* $X \in \{N, D\}$,

$$\mathcal{L}[\text{FV3-SNSP}k\text{-HXFA}^c] \subseteq \mathcal{L}[\text{CS-FV3-XFA}(2k)^c]$$

Proof. Let M be an FV3-SNSPk-HFA^c. We will construct a CS-FV3-XFA$(2k)^c$ M' to simulate M. M' acts as follws:

1. M' simulates the moves of the reading head of M and all the east,west,south,or north moves of counting heads of M by using its $(k+1)$ finite automata.
2. M' simulates all the down moves of counting heads of M by making the east moves of input heads of its other $(k-1)$ finite automata.
3. During the simulation, if M moves its reading head down, then M' makes all of input heads of finite automata of M' move down so that all the automata of M' can keep their input heads on the same plane and can communicate with each other in that plane.

It is easy to see that M' can simulate M. □

Lemma 3.2. *For any* $k \geq 1$ *and any* $X \in \{N, D\}$,

$$\mathcal{L}[\text{CS-FV3-XFA}(k)^c] \subseteq \mathcal{L}[\text{FV3-SNSP}(2k^2 - k + 1)\text{-HXFA}^c].$$

Proof. Let M = (FA$_1$,FA$_2$,...,FA$_k$)be a CS-FV3-XFA$(k)^c$. We will construct an FV3-SNSP$(2k^2-k+1)$-HXFA^c M' to simulate M. Let R denote the reading head of M', and $h_1, h_2, \ldots, h_{2k^2-k}$ denote the $2k^2 - k$ counting heads of M'. M' acts as follws:

1. M' stores the internal states of FA$_1$, FA_2,...,FA$_k$ in its finite contorol.
2. For each plane of the input tape:
   - **(a)** M' simulates the east,west,south,or north moves of input heads of FA$_1$, FA_2,...,FA$_k$ by using R and $h_1, h_2, \ldots, h_k$.
   - **(b)** M' stores in its finite control the internal state of each FA$_i$, $1 \leq i \leq k$, when the input head of FA$_i$ leaves the plane and the order, $(d_1, d_2, \ldots, d_k)$, in which the input heads of FA$_1$, FA$_2$, ..., FA$_k$ leave the plane subsequently (i.e., FA$_{d_1}$ firstly moves its input head down from the plane. FA$_{d_2}$ secondly moves its input head down from the plane, and so on.), and M' keeps the position where the input head of each FA$_i$, $1 \leq i \leq k$, leaves the plane by the positions of $h_1, h_2, \ldots, h_k$.
   - **(c)** Furthermore, for each $i (1 \leq i \leq k-1)$, the interval between the times at which FA$_{d_i}$ and FA$_{d_{i+1}}$ move their input heads down from the plane is stored by a counter with $O(n^{4k})$ space bound, which can be realized by using $h_{(2i-1)k-1}$, $h_{(2i-1)k-2}, \ldots, h_{(2i-1)k}$, where n is the number of rows (or columns or planes) of the input tape.

Note that M works in $O(n^{4k})$ time, that is, if an input tape with n rows (or columns or planes) is accepted by M, then it can be accepted by M in $O(n^{4k})$ time. Thus, it is easy to verify that M' can simulate M. □

From [5],it follows that $\cup_{1 \leq k < \infty} \mathcal{L}$[FV3-SP$k$-HXFA^c] $= \cup_{1 \leq k < \infty} \mathcal{L}$[FV3-SNSP$k$-HXFA^c] for any $X \in \{N, D\}$. Combining this result with Lemmas 3.1 and 3.2, we have the follwing thorem.

Theorem 3.1. $\cup_{1 < \infty} \mathcal{L}$[FV3-SP$k$-HXFA^c] $= \cup_{1 \leq k < \infty} \mathcal{L}$[CS-FV3-XFA$(k)^c$] for any $X \in \{N, D\}$.

Corolary 3.1. *For any* $k \geq 1$, *there is no* CS-FV3-NFA(k) *that accepts the set of connected patterns.*

Remark 3.1. It is easy to see that for each $k \leq 1$, (1)three-dimensional sensing simple k head finite automata [5] are simulated by cooperating systems of $(k+1)$ three-dimensional finite automata, and (2) cooperating systems of k three-dimensional finite automata are simulated by three-dimensional sensing simple $(k+1)$ head finite automata.

Remark 3.2. It is shown in [8] that (one-dimensional) one-way simple multihead finite automata snd cooperating systems of (one-dimensional) one-way deterministic finite automata are incomparable in accepting power. From this fact, it follows that FV3-SPMHFA's and CS-FV3-DFA's are incomparable in accepting power if the input tapes are restricted to those x such that $l_3(x) > l_1(x) = l_2(x)$. We can also show that FV3-SPMHFA's are more powerful than CS-FV3-DFA's if the input tapes are restricted to those x such that $l_3(x) < l_1(x) = l_2(x)$.

4 Five-Way versus Six-Way

In this section, we investigate the difference between the accepting powers of CS-3-DFA$(k)^c$'s [CS-3-NFA$(k)^c$'s] and CS-FV3-DFA$(k)^c$'s [CS-FV3-NFA$(k)^c$'s].

Theorem 4.1. *For each* $X \in \{N, D\}$, $\mathcal{L}$ [CS-3-DFA$(1)^c$]$- \cup_{1 \leq k < \infty} \mathcal{L}$[CS-FV3-XFA$(k)^c$] $\neq \emptyset$.

Proof. Let $T_1 = \{x \in \{0,1\}^{(3)} | (\exists m \geq 2)[l_1(x) = l_2(x) = l_3(x) = m \ \& \ x[(1,1,1),(m,m,1)] = x[(1,1,2),(m,m,2)]]\}$. Clearly, $T_1 \in \mathcal{L}$[CS-3-DFA$(1)^c$]. From [5], it is easy to see that T_1 is not in $\cup_{1 \leq k < \infty} \mathcal{L}$[FV3-SP$k$-HNFA^c]. From this fact and Theorem 3.1, the theorem follows. □

From Theorem 4.1, we can get the following corollary.

Corollary 4.1. *For each* $k \geq 1$ *and* $X \in \{N, D\}$, (1)$\mathcal{L}$[CS-TR2-XFA$(k)^c$]$\subsetneq$ $\mathcal{L}$[CS-2-XFA$(k)^c$], *and* (2) $\cup_{1 \leq k < \infty} \mathcal{L}$[CS-TR2-XFA$(k)^c$] $\subsetneq \cup_{1 \leq k < \infty} \mathcal{L}$[CS-2-XFA$(k)^c$].

5 Nondeterminism versus Determinism

In this section, we investigate the difference between the accpting powers of CS-FV3-NFA$(k)^c$'s and CS-FV3-DFA$(k)^c$'s.

Theorem 5.1. $\mathcal{L}$[CS-FV3-NFA$(1)^c$] $- \cup_{1 \leq k < \infty} \mathcal{L}$[CS-FV3-DFA$(k)^c$] $\neq \emptyset$.

Proof. Let $T_2 = \{x \in \{0,1\}^{(3)} | (\exists m \geq 2)[l_1(x) = l_2(x) = l_3(x) = m] \ \& \ \exists_i, \exists_j (1 \leq i \leq m, 1 \leq j \leq m)[x(i,j,1) = x(i,j,2) = 1]$. Clearly, $T_2 \in$ [CS-FV3-NFA$(1)^c$]. From [5], it is easy to see that T_2 is not in $\cup_{1 \leq k < \infty} \mathcal{L}$[FV3-SP$k$-HDFA^c]. From this fact and Theorem 3.1, the theorem follows. □

From Theorem 5.1, we get the following corollary.

Corollary 5.1. *For each* $k \geq 1$, (1)$\mathcal{L}$[CS-FV3-DFA$(k)^c$] $\subsetneq \mathcal{L}$[CS-FV3-NFA$(k)^c$], *and*(2) $\cup_{1 \leq k < \infty} \mathcal{L}$[CS-FV3-DFA$(k)^c$] $\subsetneq \cup_{1 \leq k < \infty} \mathcal{L}$[CS-FV3-NFA$(k)^c$].

6 Conclusion

We conclude this paper by giving several open problems except the open problem stated in the previous section.

In this paper, we introduced a cooperating system of three-dimensional finite automata, and investigated several basic accepting powers. We conclude this paper by giving an open problem as follows.

For each $k \geq 2$,

$$\mathcal{L}[\text{CS-3-DFA}(k)^c] \subsetneq \mathcal{L}[\text{CS-3-NFA}(k)^c] \ ?$$

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